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The Low-Level Bounded Model Checker LLBMC

A Precise Memory Model for LLBMC

Carsten Sinz Stephan Falke Florian Merz | October 7, 2010

VERIFICATION MEETS ALGORITHM ENGINEERING

Motivation

Buffer overflows are still the number one issue as reported in OS vendor advisories. (...) Integer overflows, barely in the top ten overall in the past few years, are number two for OS vendor advisories (in 2006), behind buffer overflows

Use-after-free vulnerability in Microsoft Internet Explorer (...) allows remote attackers to execute arbitrary code by accessing a pointer associated with a deleted object (...)

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Use-after-free vulnerability in Microsoft Internet Explorer (. . .) allows remote attackers to execute arbitrary code by accessing a pointer associated with a deleted object (. . .)

What is LLBMC?

- LLBMC = Low-Level (Software) Bounded Model Checking
  - Low-Level: Not operating on source code but on “abstract assembler”
  - Software: Programs written in C/C++/Objective C and compiled into “abstract assembler”
  - Bounded: restricted number of nested function calls and loop iterations
  - Model Checking: bit-precise static analysis
- Properties checked:
  - Built-in properties: invalid memory accesses, use-after-free, double free, range overflow, division by zero, . . .
  - User-supplied properties: `assert` statements
- Focus on memory properties

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Software Bounded Model Checking

- Programs typically deal with unbounded data structures such as linked lists, trees, etc.
- Property checking is undecidable for these programs
- Bugs manifest themselves in (typically short) finite runs of the program
- Software bounded model checking:
  - Analyze only bounded program runs
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Specifying and Verifying Properties

- Properties are formalized using `assume` and `assert` statements
  - `assume` states a pre-condition that is assumed to hold at its location
  - `assert` states a post-condition that is to be checked at its location
- The program `Prog` is correct if

$$\mathtt{Prog} \wedge \bigwedge \mathtt{assume} \Rightarrow \bigwedge \mathtt{assert}$$

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Low Level Bounded Model Checking

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- Fully supporting real-life programming languages is cumbersome
- Particularly true for C/C++/Objective C due to their complex (sometimes ambiguous) semantics
- Key idea: Do not operate on the source code directly, use a compiler intermediate language ("abstract assembler") instead
  - Well-defined, simple semantics makes logical encoding easier
  - Closer to the code that is actually run
  - Compiler optimizations etc. come "for free"
- LLBMC uses the LLVM intermediate language and compiler infrastructure
- After the logical encoding, LLBMC uses the SMT solver `Boolector` (theory of bitvectors and arrays)

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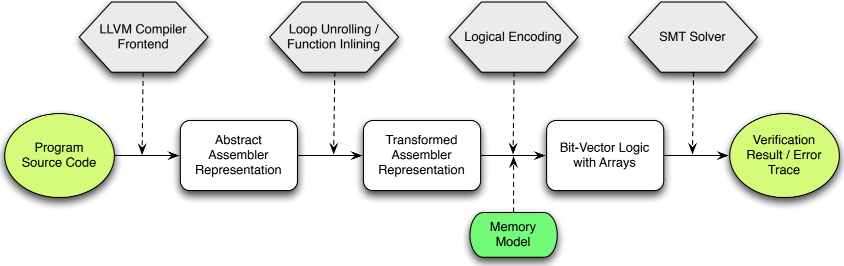
Overview of the LLBMC Approach



Memory model captures the semantics of memory accesses

Example

```c
struct S {
    int x;
    struct S *n;
};

int main(int argc, char *argv[]) {
    struct S *p, *q;

    p = malloc(sizeof(struct S));
    p->x = 5;
    p->n = NULL;

    if (argc > 1) {
        q = malloc(sizeof(struct S));
        q->x = 5;
        q->n = p;
    } else {
        q = p;
    }

    __llbmc_assert(p->x + q->x == 10);

    free(q);
    free(p);

    return 0;
}
```

```llvm
%struct.S = type { i32, %struct.S* }

define i32 @main(i32 %argc, i8** %argv) {
entry:
  %0 = call i8* @malloc(i32 8)
  %p = bitcast i8* %0 to %struct.S*
  %p.x = getelementptr %struct.S* %p, i32 0, i32 0
  store i32 5, i32* %p.x
  %p.n = getelementptr %struct.S* %p, i32 0, i32 1
  store %struct.S* null, %struct.S** %p.n
  %c.1 = icmp sgt i32 %argc, 1
  br i1 %c.1, label %if.then, label %if.end

if.then:
  %1 = call i8* @malloc(i32 8)
  %q = bitcast i8* %1 to %struct.S*
  %q.x = getelementptr %struct.S* %q, i32 0, i32 0
  store i32 5, i32* %q.x
  %q.n = getelementptr %struct.S* %q, i32 0, i32 1
  store %struct.S* %p, %struct.S** %q.n
  br label %if.end

if.end:
  %q.0 = phi %struct.S* [ %q, %if.then ], [ %p, %entry ]
  %q.0.x = getelementptr %struct.S* %q.0, i32 0, i32 0
  %2 = load i32* %p.x
  %3 = load i32* %q.0.x
  %4 = add i32 %2, %3
  %c.2 = icmp eq i32 %4, 10
  %5 = zext i1 %c.2 to i32
  call void @__llbmc_assert(i32 %5)
  %6 = bitcast %struct.S* %q.0 to i8*
  call void @free(i8* %6)
  %7 = bitcast %struct.S* %p to i8*
  call void @free(i8* %7)
  ret i32 0
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```

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  call void @free(i8* %7)
  ret i32 0
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Encoding of phi-Instructions

- The abstract assembler contains phi-instructions of the form

$$i' = \texttt{phi}[i_1, bb_1], \ldots, [i_n, bb_n]$$

where $bb_1, \ldots, bb_n$ are basic blocks

- For the logical encoding, bb_j is replaced by

$$c_{\texttt{exec}}(bb_j) \wedge t(bb_j, b)$$

where

  - $c_{\texttt{exec}}(bb_j)$ is bb_j's execution condition
  - b is the basic block containing the phi-instruction
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Elimination of branches

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- The memory can be modelled as an array of bytes
- SSA form for the memory by introducing an abstract type `memstate`:
  - Memory is accessed using `read`-instructions
  - Memory is changed using `write`-, `malloc`-, and `free`-instructions
  - `phi`-instructions for memory states are introduced
- With the encoding of `phi`-instructions and the conversion of the memory to SSA form branches can be eliminated

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- With the encoding of `phi`-instructions and the conversion of the memory to SSA form branches can be eliminated

Example

```
%struct.S = type { i32, %struct.S* }

define i32 @main(i32 %argc, i8** %argv) {
entry:
  %0 = call i8* @malloc(i32 8)
  %p = bitcast i8* %0 to %struct.S*
  %p.x = getelementptr %struct.S* %p, i32 0, i32 0
  store i32 5, i32* %p.x
  %p.n = getelementptr %struct.S* %p, i32 0, i32 1
  store %struct.S* null, %struct.S** %p.n
  %c.1 = icmp sgt i32 %argc, 1
  br i1 %c.1, label %if.then, label %if.end

if.then:
  %1 = call i8* @malloc(i32 8)
  %q = bitcast i8* %1 to %struct.S*
  %q.x = getelementptr %struct.S* %q, i32 0, i32 0
  store i32 5, i32* %q.x
  %q.n = getelementptr %struct.S* %q, i32 0, i32 1
  store %struct.S* %p, %struct.S** %q.n
  br label %if.end

if.end:
  %q.0 = phi %struct.S* [ %q, %if.then ], [ %p, %entry ]
  %q.0.x = getelementptr %struct.S* %q.0, i32 0, i32 0
  %2 = load i32* %p.x
  %3 = load i32* %q.0.x
  %4 = add i32 %2, %3
  %c.2 = icmp eq i32 %4, 10
  %5 = zext i1 %c.2 to i32
  call void @__llbmc_assert(i32 %5)
  %6 = bitcast %struct.S* %q.0 to i8*
  call void @free(i8* %6)
  %7 = bitcast %struct.S* %p to i8*
  call void @free(i8* %7)
  ret i32 0
}
```

```
struct.S = struct { i32, struct.S* }

memstate %mem0
i8* %0
memstate %mem1 = malloc(%mem0, %0, 8)
struct.S* %p = bitcast(%0)
i32* %p.x = getelementptr(%p, 0, 0)
memstate %mem2 = store(%mem1, %p.x, 5)
struct.S** %p.n = getelementptr(%p, 0, 1)
memstate %mem3 = store(%mem2, %p.n, null)
i32 %argc
i1 %c.1 = %argc > 1

i8* %1
memstate %mem4 = malloc(%mem3, %1, 8)
struct.S* %q = bitcast(%1)
i32* %q.x = getelementptr(%q, 0, 0)
memstate %mem5 = store(%mem4, %q.x, 5)
struct.S** %q.n = getelementptr(%q, 0, 1)
memstate %mem6 = store(%mem5, %q.n, %p)

memstate %mem7 = phi([%mem3, !%c.1], [%mem6, %c.1])
struct.S* %q.0 = phi([%p, !%c.1], [%q, %c.1])
i32* %q.0.x = getelementptr(%q.0, 0, 0)
i32 %2 = load(%mem7, %p.x)
i32 %3 = load(%mem7, %q.0.x)
i32 %4 = add(%2, %3)
i1 %c.2 = %4 == 10
assert(%c.2)
memstate %mem8 = free(%mem7, %q.0)
memstate %mem9 = free(%mem8, %p);
```

Example

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%struct.S = type { i32, %struct.S* }

define i32 @main(i32 %argc, i8** %argv) {
entry:
  %0 = call i8* @malloc(i32 8)
  %p = bitcast i8* %0 to %struct.S*
  %p.x = getelementptr %struct.S* %p, i32 0, i32 0
  store i32 5, i32* %p.x
  %p.n = getelementptr %struct.S* %p, i32 0, i32 1
  store %struct.S* null, %struct.S** %p.n
  %c.1 = icmp sgt i32 %argc, 1
  br i1 %c.1, label %if.then, label %if.end

if.then:
  %1 = call i8* @malloc(i32 8)
  %q = bitcast i8* %1 to %struct.S*
  %q.x = getelementptr %struct.S* %q, i32 0, i32 0
  store i32 5, i32* %q.x
  %q.n = getelementptr %struct.S* %q, i32 0, i32 1
  store %struct.S* %p, %struct.S** %q.n
  br label %if.end

if.end:
  %q.0 = phi %struct.S* [ %q, %if.then ], [ %p, %entry ]
  %q.0.x = getelementptr %struct.S* %q.0, i32 0, i32 0
  %2 = load i32* %p.x
  %3 = load i32* %q.0.x
  %4 = add i32 %2, %3
  %c.2 = icmp eq i32 %4, 10
  %5 = zext i1 %c.2 to i32
  call void @__llbmc_assert(i32 %5)
  %6 = bitcast %struct.S* %q.0 to i8*
  call void @free(i8* %6)
  %7 = bitcast %struct.S* %p to i8*
  call void @free(i8* %7)
  ret i32 0
}
```

```
struct.S = struct { i32, struct.S* }

memstate %mem0
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memstate %mem1 = malloc(%mem0, %0, 8)
struct.S* %p = bitcast(%0)
i32* %p.x = getelementptr(%p, 0, 0)
memstate %mem2 = store(%mem1, %p.x, 5)
struct.S** %p.n = getelementptr(%p, 0, 1)
memstate %mem3 = store(%mem2, %p.n, null)
i32 %argc
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i8* %1
memstate %mem4 = malloc(%mem3, %1, 8)
struct.S* %q = bitcast(%1)
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memstate %mem5 = store(%mem4, %q.x, 5)
struct.S** %q.n = getelementptr(%q, 0, 1)
memstate %mem6 = store(%mem5, %q.n, %p)

memstate %mem7 = phi([%mem3, !%c.1], [%mem6, %c.1])
struct.S* %q.0 = phi([%p, !%c.1], [%q, %c.1])
i32* %q.0.x = getelementptr(%q.0, 0, 0)
i32 %2 = load(%mem7, %p.x)
i32 %3 = load(%mem7, %q.0.x)
i32 %4 = add(%2, %3)
i1 %c.2 = %4 == 10
assert(%c.2)
memstate %mem8 = free(%mem7, %q.0)
memstate %mem9 = free(%mem8, %p);
```

Encoding Memory Constraints 1

- The following memory checks are built-in:
  - Valid read/writes (i.e., only to allocated memory)
  - Valid frees (i.e., `free` is only called for the beginning of a block of allocated memory)
  - No double frees (i.e., no memory block is `free`'d twice)
- Building blocks:
  - `valid_mem_access`(m, p, s): the range $p, \ldots, p+s-1$ is allocated in the memory state m
  - `deallocated`(m, m', p): the block beginning at p is `free`'d between m and m'
  - ...

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Memory Modification Graph

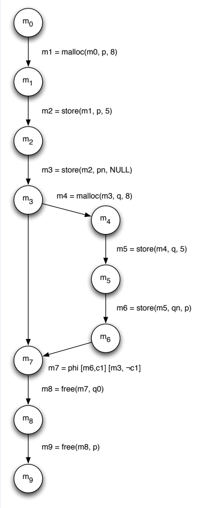
Example

```
struct.S = struct { i32, struct.S* }

memstate %mem0
i8* %0
memstate %mem1 = malloc(%mem0, %0, 8)
struct.S* %p = bitcast(%0)
i32* %p.x = getelementptr(%p, 0, 0)
memstate %mem2 = store(%mem1, %p.x, 5)
struct.S** %p.n = getelementptr(%p, 0, 1)
memstate %mem3 = store(%mem2, %p.n, null)
i32 %argc
i1 %c.1 = %argc > 1

i8* %1
memstate %mem4 = malloc(%mem3, %1, 8)
struct.S* %q = bitcast(%1)
i32* %q.x = getelementptr(%q, 0, 0)
memstate %mem5 = store(%mem4, %q.x, 5)
struct.S** %q.n = getelementptr(%q, 0, 1)
memstate %mem6 = store(%mem5, %q.n, %p)

memstate %mem7 = phi([%mem3, !%c.1], [%mem6, %c.1])
struct.S* %q.0 = phi([%p, !%c.1], [%q, %c.1])
i32* %q.0.x = getelementptr(%q.0, 0, 0)
i32 %2 = load(%mem7, %p.x)
i32 %3 = load(%mem7, %q.0.x)
i32 %4 = add(%2, %3)
i1 %c.2 = %4 == 10
assert(%c.2)
memstate %mem8 = free(%mem7, %q.0)
memstate %mem9 = free(%mem8, %p);
```



Encoding Memory Constraints 2

$m \preceq m'$: there exists a path from m to m' in the memory modification graph

$c_{\texttt{exec}}(I)$: execution condition of the (basic block containing the) instruction I

$$\texttt{deallocated}(m, m', p) \equiv \bigvee_{\substack{m \preceq m^* \preceq m' \\ I:\ m^* = \texttt{free}(\hat{m}^*, q)}} c_{\texttt{exec}}(I) \wedge p = q$$

$$\texttt{valid\_mem\_access}(m, p, s) \equiv \bigvee_{\substack{m' \preceq m \\ I:\ m' = \texttt{malloc}(\hat{m}, q, t)}} c_{\texttt{exec}}(I) \wedge (q \leq p \leq q + t - s) \wedge \neg\texttt{deallocated}(m', m, q)$$

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Encoding Memory Constraints 3

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- Each $m' = \texttt{write}(m, p, x)$ and each $x = \texttt{read}(m, p)$ is preceded by the assertion

  $$\texttt{valid\_mem\_access}(m, p, s)$$

  where s is the appropriate size
- Similar assertions are added for the other built-in memory checks
- For `malloc`-instructions, assumptions on disjointness of the allocated memory regions are added

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- Similar assertions are added for the other built-in memory checks
- For `malloc`-instructions, assumptions on disjointness of the allocated memory regions are added

Example

```
struct.S = struct { i32, struct.S* }
memstate %initialMemState
i8* %0
i1 %2 = 0x00000000 <= (void*)%0
i32 %4 = add((i32)%0, 7)
i1 %6 = 0x5fffffff >= (void*)%4
i1 %7 = (void*)%0 <= (void*)%4
i1 %8 = and(%2, %6)
i1 %9 = and(%8, %7)
assume(malloc_assume, %9, 1)
memstate %11 = malloc(heap, %initialMemState, %0, 8, 1)
i32* %p.x = getelementptr((struct.S*)%0, 0, 0)
i1 %13 = 0xbfffffff < (void*)%p.x
i32 %14 = add((i32)%p.x, 3)
i1 %16 = 0xbfffffff >= (void*)%14
i1 %17 = and(%13, %16)
i1 %18 = %0 <= %p.x
i32 %19 = add((i32)%p.x, 4)
i32 %21 = add((i32)%0, 8)
i1 %23 = (void*)%19 <= (void*)%21
i1 %24 = and(%18, %23)
i1 %25 = or(%17, %24)
assert(valid_store, %25, 1)
memstate %27 = store(%11, %p.x, 5, 1)
struct.S** %p.n = getelementptr((struct.S*)%0, 0, 1)
i1 %29 = 0xbfffffff < (void*)%p.n
i32 %30 = add((i32)%p.n, 3)
i1 %32 = 0xbfffffff >= (void*)%30
i1 %33 = and(%29, %32)
i1 %34 = %0 <= %p.n
i32 %35 = add((i32)%p.n, 4)
i1 %37 = (void*)%35 <= (void*)%21
i1 %38 = and(%34, %37)
i1 %39 = or(%33, %38)
assert(valid_store, %39, 1)
memstate %41 = store(%27, %p.n, 0x00000000, 1)
i32 %argc
i1 %c.1 = %argc > 1
i8* %42
i1 %44 = 0x00000000 <= (void*)%42
i32 %46 = add((i32)%42, 7)
i1 %48 = 0x5fffffff >= (void*)%46
i1 %49 = (void*)%42 <= (void*)%46
i1 %50 = and(%44, %48)
i1 %51 = and(%50, %49)
i32 %52 = add((i32)%42, 8)
i1 %54 = (void*)%52 <= (void*)%0
i1 %55 = (void*)%21 <= (void*)%42
i1 %56 = or(%54, %55)
i1 %57 = and(%51, %56)
assume(malloc_assume, %57, %c.1)
memstate %59 = malloc(heap, %41, %42, 8, %c.1)
i32* %q.x = getelementptr((struct.S*)%42, 0, 0)
i1 %61 = 0xbfffffff < (void*)%q.x
i32 %62 = add((i32)%q.x, 3)
i1 %64 = 0xbfffffff >= (void*)%62
i1 %65 = and(%61, %64)
i1 %66 = %0 <= %q.x
i32 %67 = add((i32)%q.x, 4)
i1 %69 = (void*)%67 <= (void*)%21
i1 %70 = and(%66, %69)
i1 %71 = %42 <= %q.x
i1 %72 = (void*)%67 <= (void*)%52
i1 %73 = and(%71, %72)
i1 %74 = and(%c.1, %73)
i1 %75 = or(%70, %74)
i1 %76 = or(%65, %75)
assert(valid_store, %76, %c.1)
```

```
assert(valid_store, %76, %c.1)
memstate %78 = store(%59, %q.x, 5, %c.1)
struct.S** %q.n = getelementptr((struct.S*)%42, 0, 1)
i1 %80 = 0xbfffffff < (void*)%q.n
i32 %81 = add((i32)%q.n, 3)
i1 %83 = 0xbfffffff >= (void*)%81
i1 %84 = and(%80, %83)
i1 %85 = %0 <= %q.n
i32 %86 = add((i32)%q.n, 4)
i1 %88 = (void*)%86 <= (void*)%21
i1 %89 = and(%85, %88)
i1 %90 = %42 <= %q.n
i1 %91 = (void*)%86 <= (void*)%52
i1 %92 = and(%90, %91)
i1 %93 = and(%c.1, %92)
i1 %94 = or(%89, %93)
i1 %95 = or(%84, %94)
assert(valid_store, %95, %c.1)
memstate %97 = store(%78, %q.n, (struct.S*)%0, %c.1)
void* %stacktopptr0 = phi([0xbfffffff, !%c.1], [0xbfffffff, %c.1])
memstate %if.end__mem = phi([%41, !%c.1], [%97, %c.1])
struct.S* %q.0 = phi([(struct.S*)%0, !%c.1], [(struct.S*)%42, %c.1])
i32* %q.0.x = getelementptr(%q.0, 0, 0)
i1 %98 = %stacktopptr0 < (void*)%p.x
i1 %99 = and(%98, %16)
i1 %100 = %42 <= %p.x
i1 %101 = (void*)%19 <= (void*)%52
i1 %102 = and(%100, %101)
i1 %103 = and(%c.1, %102)
i1 %104 = or(%24, %103)
i1 %105 = or(%99, %104)
assert(valid_load, %105, 1)
i32 %107 = load(%if.end__mem, %p.x, 1)
i1 %109 = %stacktopptr0 < (void*)%q.0.x
i32 %110 = add((i32)%q.0.x, 3)
i1 %112 = 0xbfffffff >= (void*)%110
i1 %113 = and(%109, %112)
i1 %114 = %0 <= %q.0.x
i32 %115 = add((i32)%q.0.x, 4)
i1 %117 = (void*)%115 <= (void*)%21
i1 %118 = and(%114, %117)
i1 %119 = %42 <= %q.0.x
i1 %120 = (void*)%115 <= (void*)%52
i1 %121 = and(%119, %120)
i1 %122 = and(%c.1, %121)
i1 %123 = or(%118, %122)
i1 %124 = or(%113, %123)
assert(valid_load, %124, 1)
i32 %126 = load(%if.end__mem, %q.0.x, 1)
i32 %127 = add(%107, %126)
i1 %c.2 = %127 == 10
i1 %129 = (i8*)%q.0 == %0
i1 %130 = (i8*)%q.0 == %42
i1 %131 = and(%c.1, %130)
i1 %132 = or(%129, %131)
assert(valid_free, %132, %c.2)
i1 %134 = %0 == %0
i1 %135 = %0 == (i8*)%q.0
i1 %136 = and(%c.2, %135)
i1 %138 = and(%134, !%136)
i1 %139 = %0 == %42
i1 %140 = %42 == (i8*)%q.0
i1 %141 = and(%c.2, %140)
i1 %143 = and(%139, !%141)
i1 %144 = and(%c.1, %143)
i1 %145 = or(%138, %144)
assert(valid_free, %145, %c.2)
assert(custom, 0, !%c.2)
```

Example (Memory Management)

```
struct S {
    int x;
    struct S *n;
};

int main(int argc, char *argv[]) {
    struct S *p, *q;

    p = malloc(sizeof(struct S));
    p->x = 5;
    p->n = NULL;

    if (argc > 1) {
        q = malloc(sizeof(struct S));
        q->x = 5;
        q->n = p;
    } else {
        q = p;
    }

    __llbmc_assert(p->x + q->x == 10);

    free(q);
    free(p);

    return 0;
}
```

Example (Functional Correctness)

```
int npo2(int x) {
    unsigned int i;
    x--;
    for(i = 1; i < sizeof(int) * 8; i *= 2) {
        x = x | (x >> i);
    }
    return x + 1;
}

int main(int argc, char *argv[]) {
    int x = argc;

    __llbmc_assume(x > 0 && x < (INT_MAX >> 1));

    int n = npo2(x);

    __llbmc_assert(n >= x);
    __llbmc_assert(n < (x << 1));
    __llbmc_assert((n & (n - 1)) == 0);

    return 0;
}
```

Future Work

- Optimization of memory constraints
- Discharging of simple memory constraints using:
  - Rewriting
  - Restricted linear arithmetic
  - Boolean simplification
  - ...
- Dedicated SMT solver for memory properties
- Function inlining and loop unrolling on demand
- Modular verification
- Handling system calls (strings, memory copy, etc.)

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- Discharging of simple memory constraints using:
  - Rewriting
  - Restricted linear arithmetic
  - Boolean simplification
  - . . .
- Dedicated SMT solver for memory properties
- Function inlining and loop unrolling on demand
- Modular verification
- Handling system calls (strings, memory copy, etc.)

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